

Compensations and Runtime Monitoring

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Motivation

- More widespread use of SOA, dynamic service composition, long-lived transactions, system-of-systems architectures lead to greater need for handling failure as part of a system's normal behaviour.
- Catering for failure of components is becoming more important in various scenarios:
 - Systems built of separate components may not be able to trust the success of the constituent parts.
 - Components may be discovered, invoked at runtime, not knowing enough about them at compile time.
 - Invoking multiple services (for the same computation) and using the first result would require undoing the other partial transactions.

Compensation vs Reparation

Reparation

- Classical exception and error handling
- What to do to make up for the fact that the current block has failed
- Local by its very nature
- Compensation
 - What to do if something else fails later on and I have to `undo' this block
 - Requires storing dynamic context

An Example of Compensations

- The bog-standard example: A customer buying books from an on-line bookshop
- The service proceeds as follows:
 - The customer commits an order
 - The bookshop gets payment from the customer
 - The bookshop orders a courier
 - The books are identified in the warehouse
 - The books are packed
 - The books are posted to the customer
- But errors may occur at various points in the process:
 - The bookshop realizes that one of the books is not in stock
 - The credit card payment may fail
 - The customer may cancel an order while still being processed

An Example of Compensations





The Question

How can we use information about compensations in a system to support or strengthen runtime verification?

Further Motivation

- We are applying runtime verification techniques to an industrial case study system for financial transactions.
- Ixaris Systems Ltd are provide online payments solutions and virtual credit card provider.
- Their systems already include an implementation of compensable actions and rollbacks to handle longlived transactions.



Actions performed before the failure occurs have to be "undone":



Actions performed before the failure occurs have to be "undone":



an



Actions performed before the failure occurs have to be "undone":



anbn



Actions performed before the failure occurs have to be "undone":



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Actions performed before the failure occurs have to be "undone":

- Undoing actions may involve doing something other than the inverse of the forward action:
 - chargeAcct == charge(\in 1) / refund(95c)
- Compensations may be nested:
 - (...) / (destroyAcct / recreateBlacklistedAcct; ...)
- Compensations may be overridden:
 - (createAcct; chargeAcct; makeOrder) / destroyAcct; verifyClient
- Compensations may be scoped:
 - payment; { advertisment }; delivery
- Parallelizing parts of the process makes compensation handling more involved



So Many Logics, So Little Time

- There are various flavours of compensations:
 - Process calculus style (CSP, pi-calculus based)
 - Petri net based
 - Language-based approaches
 - Deontic logics

Remember the Question?

 How can we use information about compensations in a system to support or strengthen runtime verification?

Online Runtime Verification



- Instrumented code of the monitor M is added within the running system S, sharing the same address space.
- Running synchronously with the system.
- As soon as a problem is identified, mitigation may occur without the system running further.
- But, we are effectively monitoring *S* || *M*, not *S*.
- The overheads and interaction with the system are not always acceptable in an industrial setting.

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Offline Runtime Verification



- The system produces a log to run the verification code on at a later stage.
- Running asynchronously with the system.
- Verification much more acceptable and faithful since the logging code is typically much more lightweight than the monitoring code.
- But by the time a problem is identified it may already been too late.

Quasi-Online Runtime Verification



- Monitoring is identical to offline monitoring:
 - The system produces a log at runtime.
 - The monitors run concurrently but on separate address spaces.
 - May not be in sync.

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Quasi-Online Runtime Verification



But when a problem is identified:

- The monitor may stop the system, and
- Use compensations to undo the actions performed by the system in the meantime.

Quasi-Online Runtime Verification



- In this example we would want to perform a compensation for:
 - transfer; deposit; pay

Quasi-Online Runtime Verification

Compensate for actions:

- up to the point which the system reached,
- regressing back to just before or after the error.
- Compensations may be specified by either:
 - the system, or
 - as part of the properties themselves.
- The major challenge is resuming the system from the point where it was 'rewound' to.

Runtime Verification and LARVA



```
GLOBAL {
```

```
VARIABLES { Clock t; int c = 0; }
```

```
EVENTS {
```

...

```
interact() = {*.action()}
```

 $t30() = \{t@30*60\}$

```
PROPERTY users {
STATES {
BAD { inactive badlogins }
NORMAL { loggedin }
STARTING { loggedout }
```

} TRANSITIONS { loggedout -> loggedin [goodlogin\\t.reset();] loggedout -> loggedout [badlogin\\c++;] ...



Compensation Automata

- Extends automata with compensations using hierarchical automata with three structuring elements:
 - Compensation declaration to enact a compensation of an automaton.
 - Deviation to redirect compensations.
 - Scoping of compensations.



Compensation Automata



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Compensations in Financial Transaction Systems

- We are exploring the use of LARVA with compensations to monitor transactions handled by the systems built at Ixaris.
- Using compensations as specifications of expected behaviour and to specify recovery actions can already be done by translating into base LARVA.
- We are exploring the *quasi-online approach*, in which we use the compensations already built in the system. Major challenges are:
 - Resuming the system after recovery (easy on a transaction by transaction system, but tough otherwise).
 - Making sure that rollbacks induced by the monitors do not interfere with rollbacks induced by the system.

Reactivity and Compensations

- How can compensation be encoded in a reactive setting?
- If the monitor lags behind the system by at most n time units, compensation handling uses bounded space
- What if this is not the case?
 - Enable the monitor to pause the system
 - Skim over parts of the trace
 - Breaking loops and/or using additional properties (eg idempotency of certain compensations)



Conclusions and Future Directions

- The system is up and running at Ixaris

 monitoring for properties and compensations but still working on the quasi-online approach.
- Naïve implementation of compensations may induce unbounded overhead.
- Investigating adding real-time.

Auxiliary Slides



Monitoring of Financial Transaction Systems

- Life cycle
 - Frozen or reclaimed credit cards cannot be used in financial transactions.
 - The states in the life-cycle of an entity (eg. user, credit card) are correctly traversed, i.e. in the correct order.
- Real-time
 - After six months (but not before) of user inactivity, the user should be put in a dormant state.
 - After a year (but not before) of user inactivity, the user should be removed from the system.
- Access rights
 - A user must have the necessary right before loading money onto the credit card.
 - A user must have the necessary right before transferring money from a card to another.
- Amounts
 - The number of times a user loads money to a credit card should not exceed the stipulated amount for a day or a month.
 - The total sum of money loaded should not exceed the stipulated limit for a day or for a month.