Compensation-Aware Runtime Monitoring

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Motivation

- Monitoring induces an overhead
- Various approaches have been taken to mitigate this issue
- Sometimes even a small overhead is unacceptable

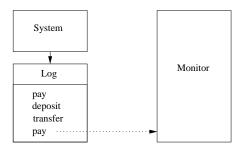
Dilemma

- Keeping system and monitor in synch
 - Slows down system
- Desynchronising the system from the monitor
 - Problems are detected late

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 - Slows down system
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 - Problems are detected late
- Synchronise only when a problem is detected

Late Detection



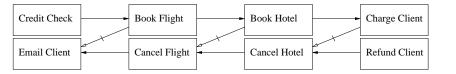
Synchronisation Technique

- In distributed games periodic synchronisation is required
- The game state of a player might have to be "reversed" to match the global state

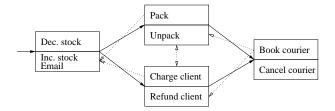
Synchronisation Technique

- In distributed games periodic synchronisation is required
- The game state of a player might have to be "reversed" to match the global state
- Compensations have been devised exactly for this purpose!

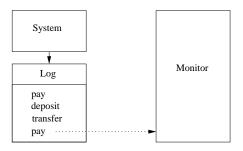
Reversing the State of a System



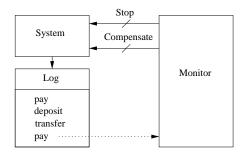
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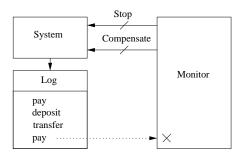
Offline Monitoring



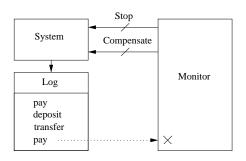
Asynchronous Monitoring with Synchronisation Capability



Asynchronous Monitoring with Synchronisation Capability



Asynchronous Monitoring with Synchronisation Capability



Upon error detection:

- Stop system
- Compensate for transfer, deposit, pay

Synchronous Monitoring

- ullet System transition system: $\sigma \stackrel{\text{a}}{\longrightarrow}_{\text{sys}} \sigma'$
- Monitor transition system: $m \xrightarrow{a}_{mon} m'$

Synchronous Monitoring

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- \bullet Stopped system state: \odot
- Broken monitor state: ⊗

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$$\text{Sync} \ \frac{\sigma \overset{\textbf{a}}{\longrightarrow}_{\textit{sys}} \ \sigma', \ m \overset{\textbf{a}}{\longrightarrow}_{\textit{mon}} \ m'}{\left(\sigma, m\right) \overset{\textbf{a}}{\longrightarrow}_{\parallel} \left(\sigma', m'\right)} \ m \neq \otimes$$

$$\text{SyncErr} \ \frac{\sigma \stackrel{a}{\longrightarrow}_{\textit{sys}} \ \sigma', \ m \stackrel{a}{\longrightarrow}_{\textit{mon}} \otimes}{(\sigma, m) \stackrel{a}{\longrightarrow}_{\parallel} (\odot, \otimes)}$$



Asynchronous Monitoring

$$\operatorname{Async}_{S} \frac{\sigma \xrightarrow{a}_{sys} \sigma'}{(\sigma, w, m) \xrightarrow{a}_{\parallel} (\sigma', wa, m)}$$

$$\operatorname{Async}_{M} \frac{m \xrightarrow{a}_{mon} m'}{(\sigma, aw, m) \xrightarrow{\tau}_{\parallel} (\sigma, w, m')}$$

$$\text{AsyncErr} \ \frac{}{\left(\sigma,w,\otimes\right) \stackrel{\tau}{\longrightarrow}_{\parallel} \left(\odot,w,\otimes\right)} \ \sigma \neq \odot$$

Asynchronous Monitoring

$$\text{Asyncs} \ \frac{\sigma \overset{\textit{a}}{\longrightarrow}_{\textit{sys}} \sigma'}{\left(\sigma, \textit{w}, \textit{m}\right) \overset{\textit{a}}{\longrightarrow}_{\parallel} \left(\sigma', \textit{wa}, \textit{m}\right)}$$

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$$\text{AsyncErr} \ \frac{}{\left(\sigma,w,\otimes\right) \stackrel{\tau}{\longrightarrow}_{\left|\right|} \left(\odot,w,\otimes\right)} \ \sigma \neq \odot$$

$$\underbrace{\text{Comp}}_{\left(\odot, wa, \otimes\right) \xrightarrow{\overline{a}}_{C} \left(\odot, w, \otimes\right)}$$



 The system and the monitor running asynchronously with compensations for synchronisation

exhibits the same behaviour as

The system running in synch with the monitor

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- abcdefxhijjīħ

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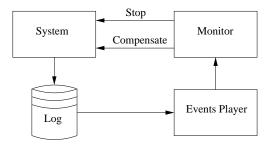
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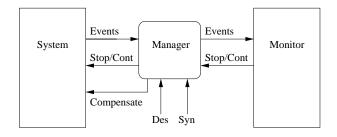
Architecture 1



Architecture 2

```
c = ok
                          :set default control to ok
while (c != stop)
 if (synch_mode)
    e = in_event()
                          ;read event from system
    c = out_event(e)
                          ;forward to monitor and get its resulting state
    out control(c)
                          ;relay control to system
 else
                          ;parallel execution
   par
      e1 = in event()
                          ;read from system
      addToBuffer(e1)
                          :store in buffer
      out control(c)
                          ;return control to system
   with
      e2 = readFromBuffer() ;read from buffer
      c = out_event(e2) ; forward to monitor and get its resulting state
end
```

Architecture 3

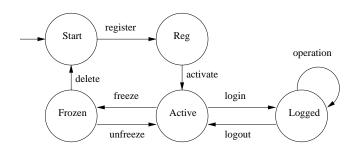


Case Study

- Life cycle properties
- Real-time properties
- Rights
- Amounts

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Conclusions

- The use of compensations to synchronise a system with its monitor
- A theory showing that monitoring asynchronously with compensation synchronisation produces the same observed behaviour as synchronous monitoring
- An architecture supporting monitor desynchronisation and resynchronisation
- The application of the architecture to an industrial case study

Future Directions

- Heuristics for deciding when to desynchronise and resynchronise
- Support a more sophisticated structure of compensations rather than the simplistic approach of having a compensation for each action

Questions

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