Compiler Theory

001 - Introduction and Course Outline Sandro Spina Department of Computer Science

Books (needed during this course)

My slides are based on the three books:

- Compilers: Principles, techniques and tools.
 Aho, Lam, Sethi, Ullman
- Engineering a Compiler, 2nd Edition.
 Keith Cooper and Linda Torczon
- Modern compiler implementation in Java.
 Andrew Appel
 - Andrew Appel
- but there are other very good books on compiler theory
- As always, the Internet is another great source of information ...

Course Sequence

- Compiling Theory ... (with practice)
 - Front end
 - Back end
- Assignment You shall be building parts of a compiler (it is very important that you have good programming skills)

Computing

- Computers are everywhere ... not just in server rooms and offices !!!
- On automobiles, telephones (mobiles), televisions, musical instruments, traffic lights, etc.
- The software that runs on them provides services such as communication, security, entertainment (eg gaming), amongst others.
- Fundumentally everything is based on the theory of computation.

Heterogeneous Computing

- **Computing hardware has evolved along many directions.**
- Today we speak about multi-core micro processors, traditional CPUs, FPGAs (reconfigurable hardware), GPUs.
- Application software can either be
 - *Control intensive (eg searching, parsing)*
 - Data intensive (eg image processing, data mining)
 - A mix of both !! (most often)
- Nowadays software needs to execute across a range of hardware devices.

What is a compiler?

- □ It is a language processor !!
- It is a <u>program</u> that can read a <u>program</u> in one language (Java, C, Lisp, C#, Pascal, C for CUDA, OpenCL, GLSL, HLSL, etc.) – the source language – and translate it into an **equivalent** <u>program</u> in another language – target language.
- Also a compiler needs to report any errors in the source program that it detects during the translation process. (e.g. a missing semicolon at the end of a statement)

Brief compiler history (based on Cooper)

- □ Appeared first in the 1950s for FORTRAN.
- □ *Early* 1960s Algol 60
- Complex machine instruction sets were created in order to mimic closely programming language features. CISC architectures.
- RISC architectures (1980s) shifted this balance towards the compiler.
- Both architectures are still widely used today ... E.g. Ipad is based on RISC and Intel's x86 on CISC.

Compiler Fundamental Principle (Cooper)



- □ An interpreter is also a language processor
- However instead of producing a target program, it directly executes the operations specified in the source program on inputs supplied by the user.
- □ A typical example of an interpreter is the JVM.
 - Java compiler translates Java source code into byte code ... which is then interpreted by a virtual machine (JIT compilers which can during runtime compile to machine code directly)

Two blocks !!

- Clearly a compiler **must** understand the syntax (valid forms) and semantics (meaning) of the input language.
- It also needs to understand the rules, syntax and meaning of the output language.



Stages of a compiler (in brief) ...

Source Program (Character Stream)
 Lexical Analysis

- Tokens (Token Stream)
 - Parser, Syntax Analyzer
- Abstract Syntax (Syntax Tree)
 - Semantic Analysis Symbol Table

Stages of a compiler (in brief) ...

- Syntax Tree
 - Intermediate Code Generator
- Intermediate Representation
 - Machine independent Optimisations
- Intermediate Representation
 - Code Generator > Target Machine Code

Eg from Aho showing compiler phases

- IMAGE FROM AHO Pg5/7
- Expression (in source code)
 *position = initial + rate * 60*

Lexical Analysis (scanning)

- □ The 1st phase of the compiler is lexical analysis (scanning)
- It reads a stream of characters (source program) and groups the characters into meaningful sequences called <u>lexemes</u>.
- For each lexeme, the lexical analyzer produces output tokens of the form
 - < token-name, attribute-value>
 - Eg <id,1> <=> <id,2> <+> <id,3> <*> <60>
- id indicates to the syntax analysis phase that we have an identifier (variable) and it's value it located in symbol table at position 1.

Syntax Analysis (parsing)

- The parser uses the first component of the tokens produced by the lexical analyzer to create a tree-like intermediate representation that depicts the grammatical structure of the token stream.
- Syntax trees are usually used where each interior node represents an operation and the children of the node represent the arguments of the operation.

Semantic Analysis

- The semantic analyzer uses the <u>syntax tree</u> and the information in the <u>symbol table</u> to check the source program for semantic consistency with the language definition.
- □ <u>Type Checking</u> !!
- Type conversions (<u>coercions</u>) ... eg if operator is applied to int and float, the compiler may convert (coerce) the integer into a floatingpoint.

Intermediate Code Generation

- Many compilers generate an explicit low-level (machine-like) intermediate representation.
- Properties
 - Easy to produce
 - Easy to translate into target machine
- Eg. Three-address code (linear representation)
 - < X = y **op** z > format
 - T2 = id3 * t1
 - T3 = id2 + t2
- We'll see how this step is useful for code optimization

Code Optimization (i)

- This stage of compilation attempts to improve the intermediate code so that better target code is generated.
- Better = faster, smaller !!
- But could also mean that target consumes less power (eg mobile phone, pda apps)
- We shall eventually look at machine-dependent and machine-independent optimizations in some detail.

Code Optimization (ii)

- Compiler optimization must meet the following design objectives:
 - It must be <u>correct</u>, that is, preserve the meaning of the compiled program.
 - It must improve the performance of <u>many</u> programs (not just a few !!)
 - Compilation <u>time</u> must be kept reasonable, and
 - Engineering effort required must be manageable

Code Generation

- The code generator takes as input an intermediate representation of the source program and maps it into the target language.
- If target language is machine code, registers or memory locations are selected for each of the variables ... the assignment of registers to hold variables is very important!
- Output might look like :
 - *LDF R2*, *id3*
 - MULF R2, R2, #60.0 (# means 60.0 is a constant)
 - LDF R1, id2
 - ADDF R1, R1, R2

Instruction Scheduling

- An important task during code generation is instruction scheduling.
- This is because the execution time of different operations can vary according to the target machine's specific performance constraints.
- Clearly, an operation cannot begin to execute until its operands are ready ... But can start multiple new operations which do not need the result of currently executing operations.
- The intuition here is that by re-ordering certain operations, the compiler minimizes the number of cycles wasted waiting for operands ...
- E.g. [load r1, add r1, load r2, mult r2, load r3, add r3] can be reordered to [load r1, load r2, load r3, add r1, mult r2, add r3]

So far we have seen (very very briefly) the phases of a compiler !!

In the next few slides we shall look at important concepts you should be aware of with respect to programming languages + compilers

Some important concepts (ii) – Symbol Table

- Compilers need to record the variable names used in the source program together with various attributes of each name (eg type, storage allocated, scope)
- Procedure names are also stored together with attributes such as number of parameters, type of arguments and method of passing each argument (by value or by reference)
- It is 'clearly' very important that the data structure which stores this information is efficient in terms of retrieval and storage of data.

Concepts (iii) – Parameter Passing

- All programming languages have a notion of a procedure (in Java we have methods, functions in C)
- Most programming languages use either call by value, call by reference or both
- <u>Call by Value</u> :- the actual parameter is evaluated (if it is an expression) or copied (if it is a variable). Both C and Java use call by value however with C we can pass a pointer to a variable and with Java many variables are really references (pointers) to arrays, strings and objects !!

Concepts (iv) – Parameter Passing

- All programming languages have a notion of a procedure (in Java we have methods, functions in C)
- Call by reference is usually an option in many programming languages (call by value is much more used)
- Call by Reference :- Changes to the formal parameter appear as changes in the actual parameter !! Eg Pascal incr (var x : int)

Concepts (v) – Memory Hierarchies

- A memory hierarchy consists of several levels of storage with different speeds and sizes, with the levels closest to the processor (registers, caches) being the fastest but smallest.
- A processor usually has a small number of registers consisting of hundreds of bytes, several levels of caches containing kilobytes to megabytes, physical memory containing megabytes to gigabytes and beyond.
- Using registers effectively is probably the single most important problem in optimizing a program.
- The compiler can improve the effectiveness of the memory hierarchy by changing the layout of the data.

Concepts (vi) – Software Tools

Lexical Analysis – LeX, FLeX, JLeX

- □ Syntax Anaysis JavaCC, SableCC
- Semantic Analysis JavaCC, SableCC
- MiniJava programming language (Appel Book)

Check Appel Pg 7

- Made up of simple statements and expressions (no loops or if stmts)
- Bottom up:
 - Binop,
 - Exp,
 - ExpList,
 - Stm.

An embedded compiler !!

- We know what a compiler is.
- Embedding with respect to language ... for example one can easily embed a parser in Haskell or Java.
- A compiler would `slightly' be more hairy to embed because you would need to generate executable at runtime using it.
- We usually think of compile-time and run-time are two separate stages but really one can be compiling at runtime !!!
- Let's take for example Janino ... (next slide)

Janino - The embedded Java compiler ... eg

ExpressionEvaluator ee = new ExpressionEvaluator(

```
"c > d ? c : d",
                           // expression
     int.class,
                           // expressionType
     new String[] { "c", "d" }, // parameterNames
     new Class[] { int.class, int.class } // parameterTypes
);
// Evaluate it with varying parameter values; very fast.
Integer res = (Integer) ee.evaluate(
     new Object[] {
                              // parameterValues
          new Integer(10),
          new Integer(11),
     }
);
System.out.println("res = " + res);
```

- Janino is used in Sunflow in order to compile shaders at runtime.
- Different shaders describing different reflection models can be loaded at run-time ... then compiled using Janino ... and executed in the same memory space of the currently executing VM.
- DirectX and OpenGL also have compilers which compile at runtime vertex and pixel shaders before loading them on the GPU.
- The point to understand here is that compilers have an extremely vast range of applicability. Not just the traditional source -> compile -> execute trace.
- Although we'll be looking mainly at those because the main concepts remain the same.

Compiler Construction Topics

- Compiler construction is a complex task !!! It combines together many aspects of Computer Science
 - Formal language theory
 - Artificial intelligence (greedy algorithms and heuristic techniques)
 - Computer architecture
- We shall cover many aspects, but will focus mostly on the front end.
- At the end of this course you will be able to take a formal specification of a language (using e.g. EBNF) and be able to use tools (compiler compilers) to generate a source-to-source compiler.

Next ...

Chapter 2 from Dragon Book

- A simple syntax-directed translator will be used to briefly introduce
 - Lexical Analysis
 - Syntax Analysis
 - Parsing
 - Intermediate Code